Parallel Computing 平行計算

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Lecture 4

Programming Shared-Address Space with OpenMP

Lecture 4

Topics

- Introduction to OpenMP
- OpenMP directives
 - specifying concurrency
 - parallel regions
 - loops, task parallelism
- Synchronization directives
 - reductions, barrier, critical, ordered
- Data handling clauses
 - shared, private, firstprivate, lastprivate
- Library primitives
- Environment variables
- Example of SPMD style OpenMP program

Introduction to OpenMP

- Open specifications for Multi Processing
- An API for explicit multi-threaded, shared memory parallelism
- Three components
 - compiler directives
 - runtime library routines
 - environment variables
- Higher-level programming model than Pthreads
 - support for concurrency, synchronization, and data handling
 - not mutexes, condition variables, data scope, and initialization
- Portable
 - API is specified for C/C++ and Fortran
 - implementations on many platforms (most Unix, Windows NT)
- Standardized

Introduction to OpenMP

• Parallelism is explicit

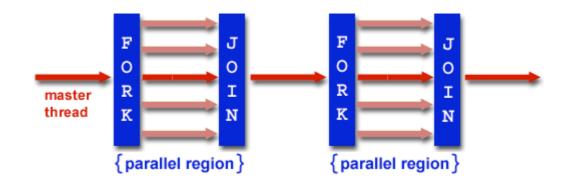
- It is not an automatic programming programming model
- programmer full control (and responsibility) over parallelization
- No data locality control
 - No guaranteed to make the most efficient use of shared memory
- Not necessarily implemented identically by all vendors
- designed for shared address spaced machines
 - Not for distributed memory parallel systems (by itself)

Introduction to OpenMP

- Advantages
 - Ease of use
 - Enables incremental parallelization of a serial program
 - Supports both coarse-grain and fine-grain parallelism
 - Portable
 - Standard

OpenMP: Fork-Join Parallelism

- OpenMP program begins execution as a single master thread
- Master thread executes sequentially until 1st parallel region
- When a parallel region is encountered, master thread
 - creates a group of threads
 - becomes the master of this group of threads
 - is assigned the thread id 0 within the group



OpenMP Directive Format

- C and C++ use compiler directives
 - prefix: #pragma ...
- Fortran uses significant comments – prefixes: !\$OMP, C\$OMP, *\$OMP
- A directive consists of a directive name followed by clauses
- C:
 - #pragma omp parallel default(shared) private(i,j)
- Fortran:
 - !\$OMP PARALLEL DEFAULT(SHARED) PRIVATE(i,j)

OpenMP parallel Region Directives

#pragma omp parallel [clause list]

Possible clauses in [clause list]

- Conditional parallelization
 - if (scalar expression)
 - determines whether the parallel construct creates threads
- Degree of concurrency
 - num_threads(integer expression)
 - Specifies the number of threads to create
- Data Handling
 - private (variable list)
 - specifies variables local to each thread
 - firstprivate (variable list)
 - similar to the private
 - private variables are initialized to variable value before the parallel directive
 - shared (variable list)
 - specifies that variables are shared across all the threads

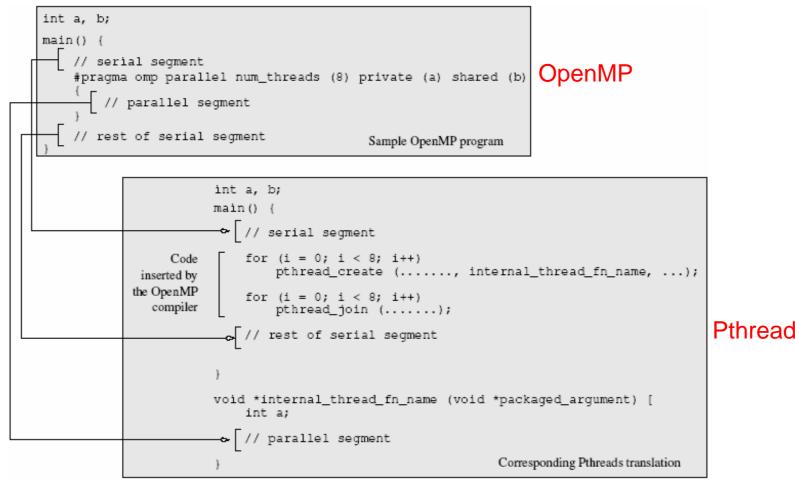
Interpreting an OpenMP Parallel Directive

```
{
   /* structured block */
}
```

Meaning

- if (is_parallel== 1) num_threads(8)
 - If the value of the variable is_parallel is one, create 8 threads
- private (a) shared (b)
 - each thread gets private copies of variables a and c
 - each thread shares a single copy of variable b
- firstprivate(c)
 - each private copy of c is initialized with the value of c in main thread when the parallel directive is encountered
- default(none)
 - default state of a variable is specified as none (rather than shared)
 - signals error if not all variables are specified as shared or private

OpenMP Programming Model



 A sample OpenMP program along with its Pthreads translation that might be performed by an OpenMP compiler.

Reduction Clause in OpenMP

- The *reduction* clause specifies how multiple local copies of a variable at different threads are combined into a single copy at the master when threads exit.
- The usage of the *reduction* clause is *reduction* (operator: variable list).
- The variables in the list are implicitly specified as being private to threads.
- The operator can be one of +, *, -, &, |, ^, &&, and ||.
 #pragma omp parallel reduction(+: sum) num_threads(8) {
 /* compute local sums here */

/*sum here contains sum of all local instances of sums */

OpenMP Programming: Example

```
/*
An OpenMP version of a threaded program to compute PI.
\#pragma omp parallel default(private) shared (npoints) \setminus
   reduction(+: sum) num threads(8)
   num threads = omp get num threads();
   sample points per thread = npoints / num threads;
   sum = 0;
   for (i = 0; i < \text{sample points per thread}; i++) {
      rand no x = (double) (rand r(\&seed)) / (double) ((2 << 14) - 1);
      rand no y = (double) (rand r(\&seed)) / (double) ((2 << 14) - 1);
      if (((rand no x - 0.5) * (rand no x - 0.5) +
          (rand no y - 0.5) * (rand no y - 0.5)) < 0.25)
          sum ++;
```

Worksharing DO/for Directive

- for directive partitions parallel iterations across threads
- DO is the analogous directive for Fortran
- Usage: #pragma omp for [clause list] /* for loop */
- Possible clauses in [clause list]
 - private, firstprivate, lastprivate
 - reduction
 - schedule, nowait, and ordered
- Implicit barrier at end of for loop

Using Worksharing for Directive

```
#pragma omp parallel default(private) shared (npoints) \
reduction(+: sum) num_threads(8)
{
                                         worksharing for divides work
  sum = 0;
  #pragma omp for
   for (i = 0; i < npoints; i++) {
      rand_no_x =(double)(rand_r(&seed))/(double)((2<<14)-1);</pre>
      rand_no_y =(double)(rand_r(&seed))/(double)((2<<14)-1);</pre>
      if (((rand_no_x - 0.5) * (rand_no_x - 0.5) +
        (rand_no_y - 0.5) * (rand_no_y - 0.5)) < 0.25)
        sum ++;
                                      Implicit barrier at end of loop
```

Example: Matrix Multiply

```
#pragma omp parallel for
for(i=0; i<n; i++)
for(j=0; j<n; j++) {
    c[i][j] =0.0;
    for (k=0; k<n; k++)
        c[i][j] += a[i][k]*b[k][i];
}</pre>
```

a,b,c are shared i,j,k are private

Private Variables

#pragma omp parallel for private(list)

- Compiler sets up a private copy of each variable in the list for each thread
- Our examples use OpenMP for and DO
- But these apply to other region and worksharing directives
- For compiler: thread has its own stack

```
for (i=0; i<n; i++) {
    tmp = a[i];
    a[i] = b[i];
    b[i] = tmp;
}</pre>
```

Swaps the values of a and b Loop-carried dependence on tmp Easily fixed by privatizing tmp

```
#pragma omp parallel for private(tmp)
for (i=0; i<n; i++) {
    tmp = a[i];
    a[i] = b[i];
    b[i] = tmp;
}</pre>
```

Removes dependence on tmp Would be more difficult to do in Pthreads

```
for (i=0; i<n; i++) {
    tmp[i] = a[i];
    a[i] = b[i];
    b[i] = tmp[i];
}</pre>
```

Requires sequential program change Wasteful in space, O(n) vs O(p)

F()

```
{ int tmp; /* local allocation on stack */
  for (i=0; i<n; i++) {
    tmp[i] = a[i];
    a[i] = b[i];
    b[i] = tmp[i];
  }
}</pre>
```

So, tmp is local to each thread

Firstprivate and Lastprivate

The initial and final values of private variables are unspecified

A firstprivate variable is private, and the private copies are initialized using its value before the loop

A lastprivate variable is private, and the thread executing the {sequentially last iteration/lexically last section} updates the version of the object outside the parallel region

Example: Firstprivate and Lastprivate

```
for(i=0; (i<n) && b[i]; i++)
a[i] = b[i];
for(i=0; j<n; j++)
a[j] = 1.0;
```

Sets all elements of a to the value of the corresponding element in b, up to first zero value in b Sets all further elements of a to 1.0

Example: Firstprivate and Lastprivate

#pragma omp parallel for lastprivate(i)
for(i=0; (i<n) && b[i]; i++)
 a[i] = b[i];
#pragma omp parallel for firstprivate(i)
for(i=0; j<n; j++)
 a[j] = 1.0;</pre>

Sets all elements of a to the value of the corresponding element in b, up to first zero value in b Sets all further elements of a to 1.0

Data Environment Directives

Private Firstprivate Lastprivate Reduction Threadprivate Copyin

For good performance, OpenMP code should use private variables whenever possible Reduces cache problems

However, this could waste a lot of memory Use of reductions also extremely important

Mapping Iterations to Threads

schedule clause of the for directive

- Recipe for mapping iterations to threads
- Usage: schedule(scheduling_class[, parameter]).
- Four scheduling classes
 - static: work partitioned at compile time
 - iterations statically divided into pieces of size chunk
 - statically assigned to threads
 - dynamic: work evenly partitioned at run time
 - iterations are divided into pieces of size chunk
 - chunks dynamically scheduled among the threads
 - when a thread finishes one chunk, it is dynamically assigned another
 - default chunk size is 1
 - guided: guided self-scheduling
 - chunk size is exponentially reduced with each dispatched piece of work
 - the default chunk size is 1
 - runtime:
 - scheduling decision from environment variable OMP_SCHEDULE
 - illegal to specify a chunk size for this clause.

Statically Mapping Iterations to Threads

- /* static scheduling of matrix multiplication loops */
- #pragma omp parallel default(private) \
- shared (a, b, c, dim) num_threads(4)
- #pragma omp for schedule(static)
- for (i = 0; i < dim; i++) {
- for (j = 0; j < dim; j++) {

```
• c(i,j) = 0;
```

```
• for (k = 0; k < dim; k++) {
```

```
c(i,j) += a(i, k) * b(k, j);
```

```
• }
```

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• static schedule maps iterations to threads at compile time

Avoiding Unwanted Synchronization

- Default: worksharing for loops end with an implicit barrier
- Often, less synchronization is appropriate
 - series of independent for-directives within a parallel construct
- nowait clause
 - modifies a for directive

- avoids implicit barrier at end of for

Avoiding Synchronization with nowait

```
#pragma omp parallel
{
#pragma omp for nowait
for (i = 0; i < nmax; i++)
    if (isEqual(name, current_list[i])
        processCurrentName(name);
#pragma omp for
for (i = 0; i < mmax; i++)
    if (isEqual(name, past_list[i])
        processPastName(name);
}</pre>
```

any thread can begin second loop immediately without waiting for other threads to finish first loop

Using the sections Directive

```
#pragma omp parallel
 #pragma omp sections
    #pragma omp section
        taskA();
   #pragma omp section
        taskB();
   #pragma omp section
        taskC();
```

parallel section encloses all parallel work sections: task parallelism three concurrent tasks

Synchronization Constructs in OpenMP

#pragma omp barrier #pragma omp single [clause list] structured block #pragma omp master structured block Use MASTER instead of SINGLE wherever possible **MASTER = IF-statement with no implicit BARRIER** equivalent to IF(omp_get_thread_num() == 0) {...} **SINGLE:** implemented like other worksharing constructs keeping track of which thread reached SINGLE first adds overhead

Synchronization Constructs in OpenMP

#pragma omp critical [(name)]
 structured block
#pragma omp ordered
 structured block

Similar to Pthreads mutex locks

critical section: like a named lock for loops with carried dependences

Example Using critical

```
#pragma omp parallel
#pragma omp for nowait shared(best_cost)
 for (i = 0; i < nmax; i++) {
     int my_cost;
     #pragma omp critical
      { if (best_cost <my_cost)
         best_cost = my_cost;
      } ....
}
critical ensures mutual exclusion
when accessing shared state
```

Example Using ordered

```
#pragma omp parallel
#pragma omp for nowait shared(best_cost)
  for (k = 0; k < nmax; k++) {
   . . .
   #pragma omp ordered
    \{a[k] = a[k-1] +
      ....
      . . .
  }
ordered ensures carried dependence does not cause a data race
```

OpenMP Library Functions

Processor count

int omp_get_num_procs(); /* # PE currently available */
int omp_in_parallel(); /* determine whether running in parallel */

Thread count and identity

/* max # threads for next parallel region. only call in serial region */
void omp_set_num_threads(int num_threads);
int omp_get_num_threads(); /*# threads currently active */
int omp_get_max_threads(); /* max # concurrent threads */
int omp_get_thread_num(); /* thread id */

OpenMP Environment Variables

• OMP_NUM_THREADS

- specifies the default number of threads for a parallel region
- OMP_SET_DYNAMIC
 - specifies if the number of threads can be dynamically changed
- OMP_NESTED
 - enables nested parallelism
- OMP_SCHEDULE
 - specifies scheduling of for-loops if the clause specifies runtime

OpenMP SPMD Style

- SPMD (Single Program Multiple Data)
- The same program on each CPU accessed different data

OpenMP SPMD Style

```
#include <omp.h>
main()
   long int i;
   long int A[100000];
   float B[1000000];
   float c[1000000];
   printf("omp_get_num_procs = %4d \n",omp_get_num_procs());
   printf("omp_get_max_threads = %4d \n",omp_get_max_threads());
   #pragma omp parallel
   ł
       #pragma omp for
     for (i=1; i<=1000000; i++)
       \{ A[i] = i; \}
        B[i] = A[i] *2.3:
}
   for(i=10; i<=100; i++) printf(" %7d ",A[i]);
   printf("\n");
}
```

```
#include <omp.h>
long int A[1000000];
int mystart, myend;
float B[1000000];
void mywork(int, int);
```

OpenMP SPMD Example

```
#pragma omp threadprivate(mystart, myend)
main()
```

{ long int i, iam; int N, nthreads, chunk,temp;

```
#pragma omp parallel private(iam, nthreads, chunk)
{ nthreads = omp_get_num_threads();
    iam = omp_get_thread_num();
    chunk = (N + nthreads - 1) / nthreads;
    mystart = iam * chunk + 1;
    temp = (iam+1) * chunk;
    myend = (temp <= N) ? temp. N;
    mywork(mystart, myend);
}
for(i=1; i<=100; i++) printf(" %7d ",A[i]);
printf("\n");</pre>
```

```
void mywork(int mystart, int myend)
{ int i;
    for (i=mystart; i<=myend; i++)
        { A[i] = i;
        B[i] = A[i] *2.3;
      }
}</pre>
```